

Secure Patient Authentication Scheme in the Healthcare System Using Symmetric Encryption

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Abstract

Recently, the incorporation of state-of-the-art technology such as Electronic Healthcare Records (EHRs), networks, and cloud computing has transformed the traditional healthcare system. However, security problems have arisen as a result of the integration of technology. Secure remote user authentication is a core part of the healthcare system to validate the user's identification via an unsecure communication network. Since then, several remote user authentication schemes have been presented, each with its own set of pros and limitations. As a result, security, malicious attacks and privacy concerns are considered one of the main challenges related to the healthcare system. In this paper, we propose a safe user authentication scheme for patients in the healthcare system that overcomes these flaws and confirms the security of the proposed work using scyther, a formal security tool. In the healthcare environment, our work provides an effective means to construct an environment capable of setting, registering, storing, searching, analyzing, authentication, and verifying electronic healthcare information in order to protect the information of patients. Furthermore, our suggested scheme uses symmetric encryption based on the crypto-hash function for accessing the anomaly of the patient's identity and One-Time Password (OTP). Towards the end of the study, the performance analysis results indicate a delicate balance of security and performance that is frequently lacking in previous works.

Keywords: Electronic Health Records, Malicious Attacks, Healthcare System, OTP Authentication.

I. INTRODUCTION

The Internet has become an integral aspect of modern life. With the rapid advancement of Internet technology, we can now provide any service from anywhere and at any time. Remote user authentication is becoming a crucial aspect of accessing valuable services or resources in the healthcare system, cloud applications, multi-server environments, and mobile devices. Remote user authentication is an important part of any security design. Authorization grants Identity-Based Privileges and audit trails are not transparent without authentication. Secrecy and privacy will be breached if we are unable to distinguish between authorized and unauthorized parties. Likewise, in order to access resources situated in faraway locations, each user must have the necessary access privileges. The use of a password-based authentication technique is one of the most simple and convenient protection mechanisms. The healthcare system, E-business, Database Management Systems, and Smart Card applications are some instances of password-based authentication schemes. There are two major issues with the

password process in the computer system. One example is that passwords are kept in database systems in plaintext that the database administrator may readily view. Another issue is that an attacker can impersonate a valid user by stealing the user ID and password from the password database. Individuals' e-health data is some of their most sensitive information. Privacy regulations such as the Health Insurance Portability and Accountability Act (HIPPA) and the General Data Protection Regulation (GDPR)[1]. The remote system should have the skill to authenticate the users. Otherwise, a discount could impersonate a legitimate user login to get access to the system[2]. They were intended to improve healthcare data governance; however, e-health data has frequently been violated. Furthermore, as the accessibility and usability of e-health data grows, so do the security attack vectors. Over the previous decade, 1.5 million medical devices have been affected owing to software flaws and wireless connections, and cloud computing services that store and analyze e-health data have become a target for massive e-health data. In 2019-2021, 41.4 million patient records



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were compromised, according to the Protenus Breach Barometer [1, 3-5].

As a result, security and privacy concerns must be addressed in order to prevent e-health data intrusions. To adequately research how to protect e-healthcare systems, security challenges, and solutions must be recognized as referred to previously. In more details, Electronic healthcare records (EHRs) are a developing technology that plays an important role in patient care. This is a medical service that allows patients to have medical appointments outside of hospitals, follow their health cases, review their medical tests, and do other medical tasks using their EHR system. Despite these benefits, studies in the literature highlight drawbacks associated with EHRs, which include disruptions to protecting patient privacy and security. Supporting a system to authenticate patients in an electronic health record (EHR) is a critical step for preventing various security risks from gaining access to patients' identities and data. Existing authentication schemes continue to have security flaws. Exchanging medical-related information or data between clients and servers exposes them to intrusion by attackers since they can be transmitted across an unstable network [6-17].

In this paper, we present a secure scheme for authenticating EHR patients and the Healthcare Center Server based on genuine member identities and verification codes. To provide robust security while maintaining satisfactory speed, the proposed work employs a lightweight crypto-hash function for the generation of One-Time Passwords (OTP) and symmetric key encryption (Counter mode (CTR mode)). As a result, the major purpose of this study is to conduct a strong authentication scheme to overcome issues in the previous works. Because current e-healthcare systems often comprise of numerous components (e.g., e-health data, medical devices, medical components (Ex. patients, doctors, employees), and cloud-health computing), there are several security risks faced these components, and our solutions are focused on authentication and privacy of patients represented the heart of healthcare system. In addition, OTP-based authentication appears to be one of the fastest-growing authentication technologies to prevents several malicious attacks such as MITM, Reply, Insider. The study demonstrate that the majority scheme of OTP-based authentication and key management solves the problems associated with dynamic ID and password schemes, especially in healthcare systems. The proposed scheme has good metrics like mutual authentication, user's anomaly, un-linkability. In practical side, the Scyther security proof was used to demonstrate our scheme's high security and resist malicious attacks. The proposed scheme has a good balances between the complexity of security and performance, can applied in healthcare system, and deals with real world data associated with the patient part.

The remainder of the article is arranged as follows: Section 2 describes the primitive tools. Section 3 shows the related work. Section 4 focused on the proposed scheme. Section 5 includes formal analysis utilizing the Scyther tool

as well as security analysis. Section 6 describes the performance results. Finally, section 7 denotes to the conclusion.

II. PRIMITIVE TOOLS

A. Counter mode (CTR mode)

Counter mode is also known as unsynchronized stream cipher mode, because the stream cipher is built from the block cipher. Here we provide a self-contained description. To encrypt using CTR mode, first, choose a uniform value $ctr \in \{0,1\}^n$. Then, by computing $y_i := F_k(ctr + i)$, where ctr and i are integers and addition is modulo 2^n , a pseudorandom stream is formed. So, $C_i := y_i \oplus m_i$ is the i^{th} ciphertext block, and the IV is transmitted as part of the ciphertext once more. It's worth repeating that F doesn't have to be invertible or even a permutation to be decrypted. The created stream may be trimmed to exactly the plaintext length. Preprocessing can be used to generate the pseudorandom stream before the message is known, and the stateful variation of CTR mode is secure, just like OFB mode, another "streamcipher" method. In contrast to the other safe modes covered so far, the CTR mode has the benefit of being able to fully parallelize encryption and decryption since all blocks of the pseudorandom stream may be calculated independently of one another. In contrast to OFB, it is also possible to decipher the cipher text's i^{th} block using just one F evaluation. If F is a pseudorandom function, then CTR mode is secure. CTR mode is appealing because of these qualities [18].

B. Crypto Hash Function

The process of turning a string of characters into a fixed-length value or key that represents the original string is referred to as hashing. The hash function indexes the original value or key and then accesses the data associated with that value or key. These which are used in cryptography should be called "one-way hash functions" can use to figure out the hash value for a given input. In the opposite way, it must be impossible to find an input for a given value that has the same hash value as that value (this is referred to as a one-way characteristic). The SHA-2 hash algorithm is used to encrypt data such as passwords. SHA-2 is a fantastic technique to ensure the security of your data, but it takes a long time to complete [19, 20].

C. The One-Time Pad

In 1917, Vernam invented the one-time pad, a totally private encryption technique. There was no indication that the proposal was absolutely hidden at the time Vernam presented it; in fact, no one knew what perfect secrecy meant. Shannon introduced the notion of complete secrecy and proved that the one-time pad delivers that degree of security some 25 years later. Use of the pad just once The assumption behind the OTP is that the encryption key must be at least as lengthy as the plaintext message and comprised of really random digits. Each letter of the plaintext is "added" to one element from the OTP using modulo-addition. When the key

is unknown, a cipher text is produced that has no relation to the plaintext. At the receiving end, the same OTP is used to retrieve the original plaintext [21].

III. RELATED WORKS

Several techniques have been presented to overcome the issues associated with cloud authentication and access control. Currently, the practice of medicine applications has sparked an interest among researchers. One of the most important topics addressed is the confidentiality of patient data, since the patient's data is private and confidential. Similarly, the patient may wish to conceal their genuine identity in order to preserve their privacy [4]. The authors of [22] covered numerous security needs, such as EHR storage security, malicious code prevention, protected access rights management, and other factors to secure the health information system. However, they did not propose a feasible plan for a patient to move their EHR to a health information system. We may envisage a scenario in which a hospital uses just the aforementioned simple procedures to construct its own health information system with no security mechanism. Furthermore, under this method, it is not viable for each patient to execute their own EHR exchange. In contrast, the authors of [4, 23-25] proposed that each patient's health records might be portable and saved on a flash drive. This is an interesting concept, but it is currently challenging to execute. There are several security concerns to address, such as portable device security and patient medical information access privileges. More security methods, however, are required to address these types of security vulnerabilities. Furthermore, several patient authentication strategies for e-health systems have been developed [26-30]. The systems in [27, 28] were vulnerable to a user impersonation attack and did not provide session key establishment with formal security evidence. The authors of [3, 29] did not provide a forward secrecy proof for the establishment of the session key. For cloud-assisted wearable devices, Liu et al. [31] advocated local and distant authentication. To achieve mutual authentication between wearable device and a smartphone, the local authentication protocol employs a hash-based selective disclosure method and a Chebyshev chaotic map. Following local authentication, the cloud performs remote authentication of the device using a yoking-proof algorithm.

A mutual authentication procedure was developed in a few investigations [32-34]. Based on Chebyshev chaotic maps and Diffie-Hellman key exchange, Li et al. [35] suggested a mutual authentication protocol and key agreement technique. Only approved doctors and medical personnel would have access to patients' health data acquired by body sensors in the planned medical system. A digital signature was also used to assure non-repudiation of the doctor's diagnosis. Cheng et al. [36] used blockchain to bypass a mutual authentication scheme's reliance on a trusted third party.

As a result, the primary goal of this research is to develop a robust authentication method based on cryptosystem tools

to address concerns that have arisen in the previous works. Because today's e-healthcare systems generally include several components (e.g., e-health data, medical devices, medical components, and cloud-health computing), security threats and demands differ, and our solutions are centered on patient authentication and privacy. Based on the foregoing notions, we present an improved anonymous user authentication and key agreement approach for health monitoring. In the following security analysis, we showed the security of our protocol using security analysis and the Syther tool. The results of the performance comparison and efficiency analysis reveal that the proposed scheme provides a higher level of security while preserving computational efficiency.

IV. THE PROPOSED SCHEME

In this section, we present the strong healthcare authentication scheme based on five phases: Setup, Registration, Login and Authentication, Healthcare, and key management. Our work offers the healthcare scheme permitting to six main elements: Cloud Healthcare Server (*CHS*), Key Generator Center (*KGC*), Users (U_i), Patient (P_i), Administrator (*Adm*), and Doctor (D_i). The main goal of the current scheme is offered secure environment for exchanging components' data of the proposed scheme. Additionally, this work has numerous benefits such as mutual authentication, key management, password anonymity, as well as, can resist familiar malicious attacks such as insider, MITM, Reply, Impersonate, and other. The characters used in the current work are conversed in Table 1.

TABLE 1
NOTATION USED IN THE PROPOSED SCHEME

Symbol	Description
<i>CHS</i>	Cloud Healthcare Server.
<i>KGC</i>	Key Generator Center.
U_i	The user.
P_i	Patient in the system.
<i>Adm</i>	Administrator in the system.
D_i	Doctor in the system.
\oplus	XOR operation.
MITM	Man-In the middle attack.
CTR mode	Counter mode.
EHR_i	Electronic healthcare record.
$h(.)$	One-way hash function.
ID_{P_i}	Identity of patient P_i .
PW_{P_i}	Password of patient P_i .
VC, VC'', VC'	Verification code.
SK_{P_i}	Shared key between P_i and <i>CHS</i> .
<i>OFB</i>	Output Feedback mode.

r_i	The one-time random number generated by user.
	Concatenation operation.

A. Setup Phase

The setup phase considers the first step of the presented work. The *KGC* is a trust third party created all the security parameters and depends on one-way hash function $h(\cdot)$, symmetric encryption $Enc()$ /decryption ($Dec()$). Then, *KGC* creates $(SK_{P_i} \in \mathbb{Z})$ to encrypt/ decrypt data between P_i and *CHS*; where the symmetric key encryption is Counter mode (CTR mode). This type of encryption employs the block cipher algorithm, which offers significant efficiency gains over traditional encryption options without jeopardizing security. Its fitted security, in particular, has been proven. Second, the majority of the objections leveled against CTR mode are unfounded.

B. Registration Phase

User (P_i) registers his main information in *CHS* by performing the following steps:

1. P_i Chooses his identity (ID_{P_i}) and password (PW_{P_i}) by using the main website of health care institute. Also, P_i records information about his doctor and relatives (Electronic HealthCare Record (EHR_i)). EHR_i includes phone numbers of doctors and relatives, Name of patient, Pathological case, E-mail, and others.

2. P_i computes the following anonymous parameters based on the following equations:

$$AID_{P_i} = h(ID_{P_i} || SK_{P_i}) \dots Eq(1)$$

$$APW_{P_i} = h(PW_{P_i} || SK_{P_i}) \dots Eq(2)$$

3. P_i submits $(AID_{P_i}, APW_{P_i}, HCR_i)$ to *CHS*.

4. *CHS* verifies its database to check if P_i is previously registered. If so, *CHS* terminates this phase. Otherwise, the *CHS* ds a new patient's information ($AID_{P_i}, APW_{P_i}, EHR_i$) in the main secure database.

C. Login and Authentication Phase

The patient (P_i) wishes to login system for checking his EHR_i , receiving report from his doctor or sending quires to his doctor. Therefore, it is necessary to ensure from the authority of P_i to allow him accessing to the system. The important steps that use in current phase as follows:

1. P_i enters his username (ID_{P_i}) and password (PW_{P_i}), generates integer random number $r_i \in \mathbb{Z}$, and computes an anonymity of identity and anonymity one time password $AID_{P_i} = h(ID_{P_i} || SK_{P_i})$ and $APW_{P_i} = h(h(PW_{P_i} || SK_{P_i}) || r_i)$, respectively.

2. P_i encrypts r_i using symmetric key encryption (CTR mode), $E_{P_i} = Enc_{SK_{P_i}}(r_i)$.

3. P_i sends his login request $\langle AID_{P_i}, APW_{P_i}, E_{P_i} \rangle$ to *CHS*.

4. In the cloud health server (*CHS*), he checks patient's; if he was found in the database of *CHS* or no based on AID_{P_i} for

obtaining SK_{P_i} . If false, he terminates this phase. Otherwise, *CHS* restores random number by decrypting E_{P_i} where $r'_i = Dec_{SK_{P_i}}(E_{P_i})$.

5. APW'_{P_i} and APW_{P_i} , if it is valid then *CHS* sends challenge as verification code (VC) to the E-mail of P_i .

6. At this moment, P_i restores verification code (VC'') via his E-mail and encrypts VC'' using $E'_{P_i} = Enc_{SK_{P_i}}(VC')$. Then, he sends E'_{P_i} to *CHS*.

7. *CHS* computes $VC'' = Dec_{SK_{P_i}}(E'_{P_i})$ and compares between VC'' and VC ; if true then he gives permission to P_i for entering the system and applying the main operations included healthcare services at the Healthcare phase. Otherwise, He terminate the current phase.

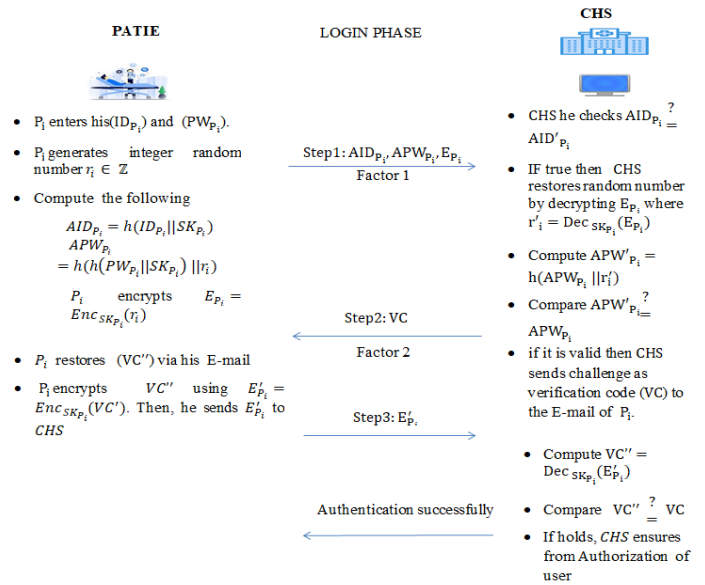


Fig.1: Login and Authentication phase

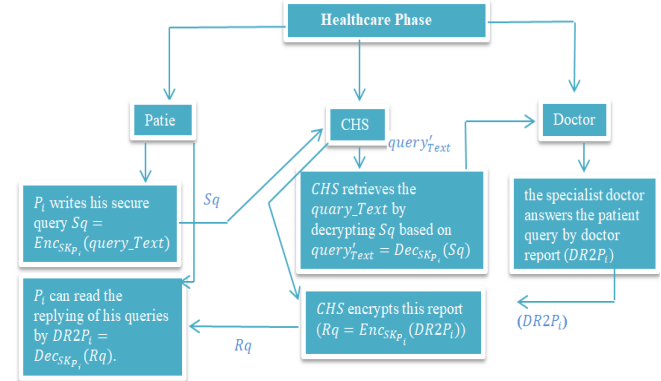


Fig.2: Flow chart of Healthcare phase for patient

D. Healthcare Phase

There are main medical services presented to each patient P_i , we can list these service in the below steps:

1. Health promotion: P_i keeps the attachment with the health foundation.

2. Disease prevention: P_i resists the disease based on devices of doctors ($D_1, D_2, D_3 \dots, D_n$) and avoids all foods caused from the severity of chronic diseases.
3. Laboratory and diagnostic care: P_i can receive the results of his test directly in his account (EHR_i). In the other side, early examination through the symptoms associated with chronic diseases by sending his symptoms to the health foundation.
4. Remote emergency and inpatient services: P_i can get full services to the case emergency to take first aid and treatment and prepare logistics before arriving at the hospital.
5. The mechanism of sending a query from P_i to CHS and vice versa. This query represents "*Inquiries about symptoms of a specific disease or other*". However, P_i writes his secure query $Sq = Enc_{SK_{P_i}}(query_Text)$ and sends Sq to CHS .
6. In the side of CHS , he retrieves the $query_Text$ by decrypting Sq based on $query'_{Text} = Dec_{SK_{P_i}}(Sq)$. Then, CHS sends the $query'_{Text}$ to the concerned department in the health foundation. After that, the specialist doctor answers the patient query by doctor report ($DR2P_i$) and then CHS encrypts this report ($Rq = Enc_{SK_{P_i}}(DR2P_i)$) and resubmits Rq to P_i .
7. Upon receiving Rq , P_i can read the replying of his queries by $DR2P_i = Dec_{SK_{P_i}}(Rq)$.

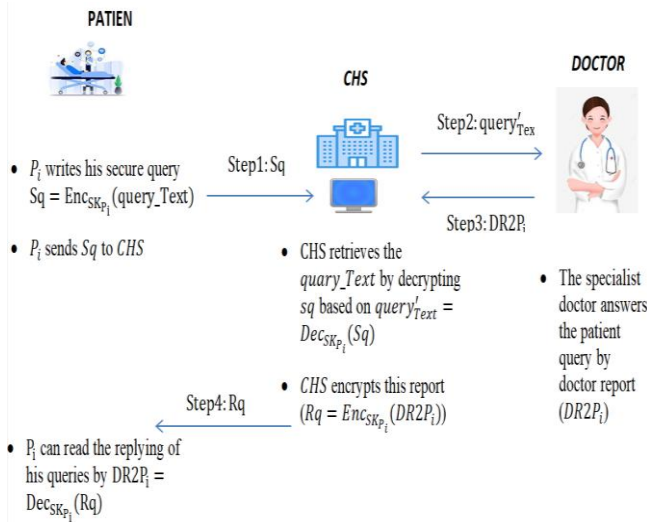


Fig.3: Healthcare phase for patient

5. Key Management Phase

Here, the main parties have an agreement to generate once key for each login request based on (SK_{P_i}, r_i) . In the moment of successful login of patient, the main parties (P_i, CHS) performs the following points to apply this phase.

1. Patient (P_i) side computes $SK_{P_i} = SK_{P_i} \oplus r_i$.
2. Cloud Healthcare Server (CHS) side computes $SK_{P_i} = SK_{P_i} \oplus r'_i$.

V. FORMAL SECURITY ANALYSIS WITH SCYTHER TOOL

Currently, we are focusing on establishing that the suggested method can withstand severe assaults such as phishing, man-in-the-middle (MITM), replay attacks, and eavesdropping attacks. Furthermore, our work incorporates a number of security features. We do the following analysis of the suggested scheme. It considers as a cryptography tool used for formal security analysis and proof to ensure from security, resisting well-known malicious attacks, correctness of the messages via communication channel applied with the perfect cryptography functions such as hash function, EMAIL, encryption, and decryption. Finally, this tool gains the proposed scheme a guarantee to implement in the safe condition and can resist the familiar attacks until the attacker is accessible. Currently, the proposed scheme has been written in SPDL language and the viewed the results in the state of (*Automatic Claim*) and (*Verification Claim*).

1. **Verification of Claim :** Scyther's input language allows security features to be specified in terms of claimed events. For example, one may argue in a role definition that a particular value is confidential (confidentiality) and that certain traits should apply to communication partners (authentication). Scyther can be used to verify or disprove these traits.

2. **Automatic of Claim:** If the protocol specification lacks security assertions, Scyther can generate them automatically. Verification claims assert that the protocol's putative communication partners must have followed it as intended at the conclusion of each role. All parameters and locally produced values are likewise subject to confidentiality claims (nonce). Scyther evaluates the expanded protocol description, in the same manner, he did previously. This allows users to quickly study a protocol's characteristics. Based on the scyther tool, our work resists harmful attacks such as MITM, Insider, Replay, Spoofing, and Impersonation. The login and authentication phases are depicted in Figure 4 [37].

We notice that the proposed scheme has protect against malicious attacks as above mention. So, SPDL supports many major crypto functions like sending or receiving messages among components, and the roles of each component. When we strip the proposed system from security factions like crypto hash, encryption, we will notice the apparent weakness of the system, as Figures 5 and 6.

Figure 5 represents the result of an safe system that does not contain perfect cryptography functions. Therefore, this system becomes insecure and vulnerable to face the malicious attacks (see Figure 6).

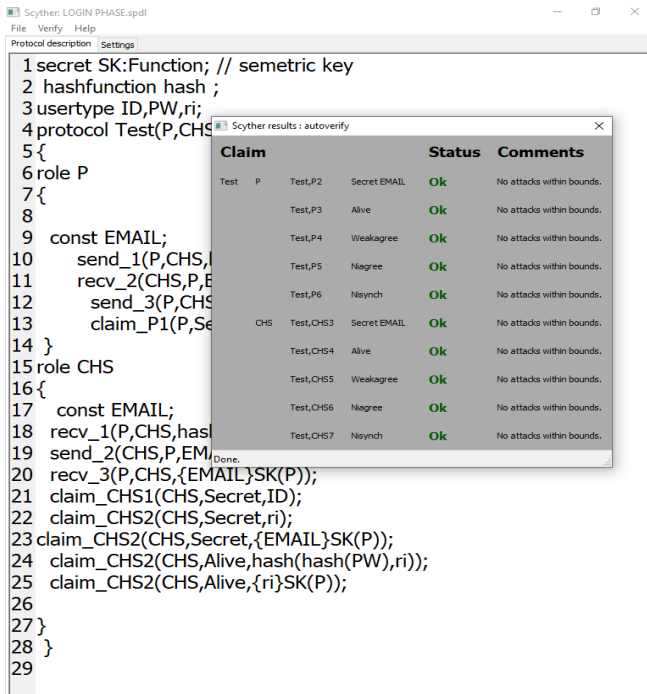


Fig.4: Login and Authentication phase that cannot be attacked

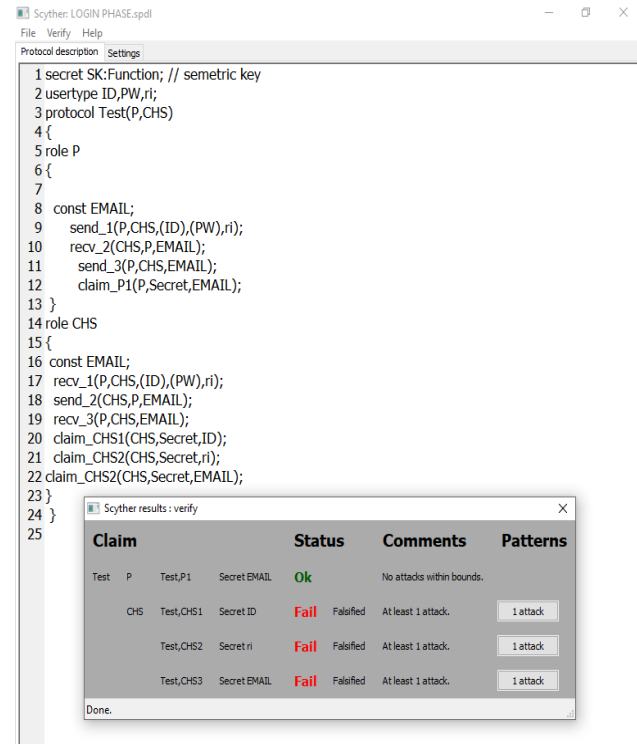


Fig. 5: Login and authentication phase that can be attacked

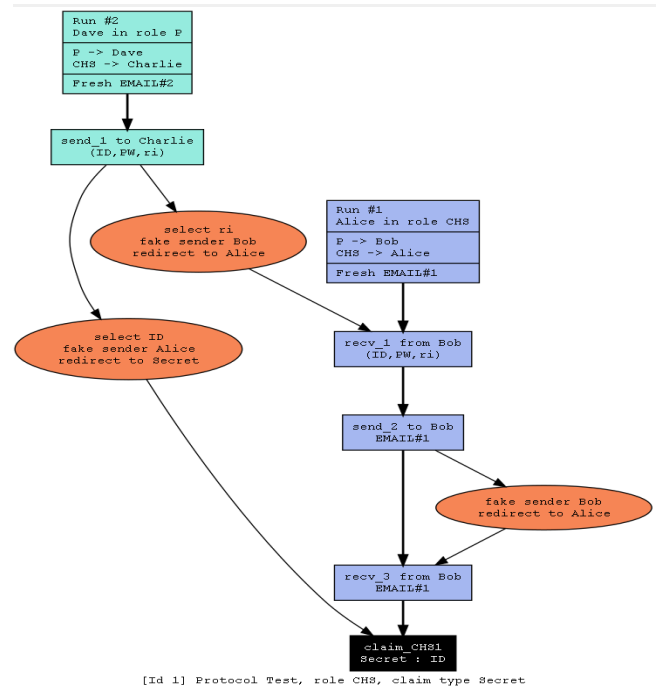


Fig. 6: Model checking of the login and authentication phase.

Figure 7 refers to the code of the login and authentication phase.

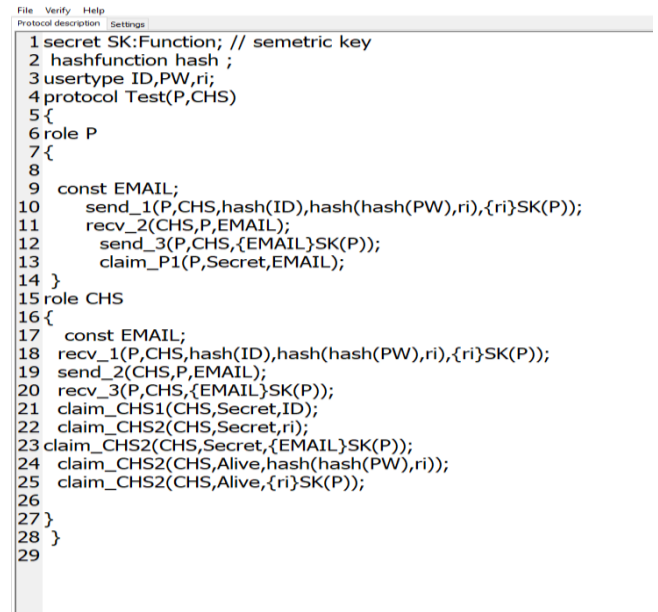


Fig. 7: The Proposed Scheme In SPDL-Scyther.

In the healthcare phase, we apply the same mechanism to check the correctness, security, verification of this phase. The figure 8 demonstrates the safety and security of the healthcare phase.

Scyther: NEWHEALTH CARE CORRECT.spdill

File Verify Help

Protocol description Settings

1 usertype DR2P;
2 secret SK:Function; // semetric key
3
4 protocol Test(P,CHS,D)
5 {
6 role P
7 {
8 secret query;
9 secret Rq;
10
11 send_1(P,CHS,{query}SK(P));
12 read_4(CHS,P,{DR2P}SK(P));
13 claim_P1(P,Secret,SK);
14 claim_P2(P,Alive,{query}SK(P));
15 claim_P3(P,Alive,{DR2P}SK(P));
16
17 }
18 role CHS
19 { secret query;
20 secret Rq;
21 read_1(P,CHS,{query}SK(P));
22 send_2(CHS,D,{query}SK(P));
23 read_3(D,CHS,Rq);
24 send_4(CHS,P,{DR2P}SK(P));
25 claim_CHS1(CHS,Secret,Rq);
26 claim_CHS2(CHS,Secret,SK);
27 claim_CHS3(CHS,Alive,{query}SK(P));
28 claim_CHS4(CHS,Alive,{DR2P}SK(P));
29 }
30 role D
31 { secret query;
32 secret Rq;
33
34 read_2(CHS,D,{query}SK(P));
35 send_3(D,CHS,Rq);
36
37 claim_D1(D,Secret,Rq);
38 claim_D2(D,Alive,{query}SK(P));
39 claim_D3(D,Secret,SK);
40 }
41 }

Scyther results : autoverify

Claim				Status	Comments
K	P	Test_P4	Secret Rq	Ok	No attacks within bounds.
		Test_P5	Secret query	Ok	No attacks within bounds.
		Test_P6	Alive	Ok	No attacks within bounds.
		Test_P7	Weakagree	Ok	No attacks within bounds.
		Test_P8	Niagree	Ok	No attacks within bounds.
		Test_P9	Nsynch	Ok	No attacks within bounds.
	CHS	Test_CHS5	Secret Rq	Ok	No attacks within bounds.
		Test_CHS6	Secret query	Ok	No attacks within bounds.
		Test_CHS7	Alive	Ok	No attacks within bounds.
D		Test_CHS8	Weakagree	Ok	No attacks within bounds.
		Test_CHS9	Niagree	Ok	No attacks within bounds.
		Test_CHS10	Nsynch	Ok	No attacks within bounds.
	D	Test_D4	Secret Rq	Ok	No attacks within bounds.
		Test_D5	Secret query	Ok	No attacks within bounds.
		Test_D6	Alive	Ok	No attacks within bounds.
		Test_D7	Weakagree	Ok	No attacks within bounds.
		Test_D8	Niagree	Ok	No attacks within bounds.
		Test_D9	Nsynch	Ok	No attacks within bounds.

Done.

Fig. 8: Healthcare phase that cannot be attacked.

Figure 9 demonstrates the result of a safe system without faultless cryptographic mechanisms. As a result, this method becomes subject to malicious assaults and is rendered unsecure (see Figure 10).

SySyn: NEW:HCARE CORRECT:Upl

File_Verify_Help

Protocol description Settings

```
1 secret SK:Function; // semetric key
```

```
2 usertype query,DR2P,Rq;
```

```
3 protocol Test(P,CHS,D)
```

```
4 {
```

```
5 role P
```

```
6 {
```

```
7 send_1(P,CHS,query);
```

```
8 read_4(CHS,P,DR2P);
```

```
9 claim_P2(P,Secret,query);
```

```
10 claim_P3(P,Secret,DR2P);
```

```
11
```

```
12 }  
13 role CHS
```

```
14 {
```

```
15 read_1(P,CHS,query);
```

```
16 send_2(CHS,D,query);
```

```
17 read_3(D,CHS,Rq);
```

```
18 send_4(CHS,P,DR2P);
```

```
19 claim_CHS1(CHS,Secret,Rq);
```

```
20 claim_CHS3(CHS,Secret,query);
```

```
21 claim_CHS4(CHS,Secret,DR2P);
```

```
22 }  
23 role D
```

```
24 {
```

```
25 read_2(CHS,D,query);
```

```
26 send_3(D,CHS,Rq);
```

```
27 claim_D1(D,Secret,Rq);
```

```
28 claim_D2(D,Secret,query);
```

```
29 }  
30 }
```

III SySyn results: verify

Claim	Status	Comments	Patterns
Test P Test,P2 Secret query	Fail	Failed At least 1 attack.	1 attack
Test,P3 Secret DR2P	Fail	Failed At least 1 attack.	1 attack
CHS Test,CHS1 Secret Rq	Fail	Failed Exactly 1 attack.	1 attack
Test,CHS3 Secret query	Fail	Failed Exactly 1 attack.	1 attack
Test,CHS4 Secret DR2P	Fail	Failed At least 1 attack.	1 attack
D Test,D1 Secret Rq	Fail	Failed At least 1 attack.	1 attack
Test,D2 Secret query	Fail	Failed Exactly 1 attack.	1 attack
Done.			

Fig. 9: Healthcare phase that can be attacked.



Fig.10: Model checking of the healthcare phase.

Protocol description	Settings
1 usertype DR2P; 2 secret SK:Function; // semetric key 3 protocol Test(P,CHS,D) 4 { 5 role P 6 { secret query; 7 secret Rq; 8 send_1(P,CHS,{query}SK(P)); 9 read_4(CHS,P,{DR2P}SK(P)); 10 claim_P1(P,Secret,SK); 11 claim_P2(P,Alive,{query}SK(P)); 12 claim_P3(P,Alive,{DR2P}SK(P)); 13 } 14 role CHS 15 { secret query; 16 secret Rq; 17 read_1(P,CHS,{query}SK(P)); 18 send_2(CHS,D,{query}SK(P)); 19 read_3(D,CHS,Rq); 20 send_4(CHS,P,{DR2P}SK(P)); 21 claim_CHS1(CHS,Secret,Rq); 22 claim_CHS2(CHS,Secret,SK); 23 claim_CHS3(CHS,Alive,{query}SK(P)); 24 claim_CHS4(CHS,Alive,{DR2P}SK(P)); 25 } 26 role D 27 secret Rq; 28 read_2(CHS,D,{query}SK(P)); 29 send_3(D,CHS,Rq); 30 claim_D1(D,Secret,Rq); 31 claim_D2(D,Alive,{query}SK(P)); 32 claim_D3(D,Secret,SK); 33 { secret query; 34 secret Rq; 35 read_2(CHS,D,{query}SK(P)); 36 send_3(D,CHS,Rq); 37 claim_D1(D,Secret,Rq); 38 claim_D2(D,Alive,{query}SK(P)); 39 claim_D3(D,Secret,SK); 40 } 41 }	

Fig. 11: The Proposed Scheme In SPDL-Scyther

VI. PERFORMANCE ANALYSIS

A. Computation Cost

The computational cost is used to determine the time complexity of the proposed scheme. Table 2 compares the computational costs of the most important related scheme and our work and Figure 12, as well as our work with other related works. Table 3 shows key security feature comparisons between the proposed scheme and previous works. Moreover, depending on [38], the processing times for the basic functions are approximately as follows. Using the rules in the following:

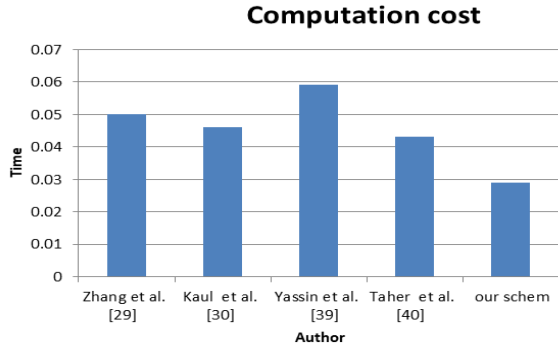


Fig. 12: Computation Cost Comparison

TABLE 2
COMPUTATION COST COMPARISON WITH OTHER RELATED WORKS

Term	Meaning	Time needed
T_h	The time allotted to the crypto hash function	0.0023 ms
T_{\oplus}	The processing time for the XOR operation	Negligible
T_{Enc}	The processing time for a symmetric encryption function.	0.0046 ms
T_{Dec}	The processing time for a symmetric decryption Function.	0.0046 ms
$T_{ }$	The processing time for the Concatenation operation.	Negligible

According to the aforementioned comparisons, the suggested system has a lower time complexity ($5T_h + 2T_{Enc} + 5T_{||} + 2T_{Dec}$) = 0.0299 than previous relevant studies. We see that the suggested system has a fair mix of performance and security aspects (see Table 3).

TABLE 3
COMPARING OF THE COMPUTATIONAL COST

Scheme	Registration Phase	Login and Authentication Phases	Total Cost
Zhang et al. [29]	$4T_h + 5T_{\oplus} + 5T_{ }$	$18T_h + 27T_{\oplus} + 19T_{ }$	$22T_h + 32T_{\oplus} + 24T_{ }$ ≈ 0.0506
Kaul et al. [30]	$6T_h + 6T_{\oplus} + 6T_{ } + 1T_{Dec} + 1T_{Enc}$	$10T_h + 20T_{\oplus} + 10T_{ }$	$16T_h + 26T_{\oplus} + 16T_{ } + 1T_{Dec} + 1T_{Enc}$ ≈ 0.046
Yassin et al. [39]	$5T_h + 2T_{\oplus} + 1T_{ }$	$13T_h + 12T_{\oplus} + 6T_{ } + 2T_{Dec} + 2T_{Enc}$	$18T_h + 14T_{\oplus} + 7T_{ } + 2T_{Dec} + 2T_{Enc}$ ≈ 0.0598
Taher et al. [40]	$5T_h + 3T_{\oplus} + 4T_{ }$	$2T_{Dec} + 2T_{Enc} + 6T_h + 3T_{\oplus} + 5T_{ }$	$11T_h + 6T_{\oplus} + 9T_{ } + 2T_{Dec} + 2T_{Enc}$ ≈ 0.0437
Our Scheme	$2T_h + 2T_{ }$	$3T_h + 3T_{ } + 2T_{Enc} + 2T_{Dec}$	$5T_h + 2T_{Enc} + 5T_{ } + 2T_{Dec}$ ≈ 0.0299

TABLE 4
COMPARISON WITH OTHER RELATED WORKS

Security Features	[4]	[22]	[29]	[30]	[32]	[34]	Our
Mutual Authentication	YES	NO	YES	YES	YES	YES	YES
Anonymous& Untraceable	NO	NO	YES	NO	NO	YES	YES
Forward Secrecy	YES	NO	YES	NO	YES	YES	YES
Key Agreement	NO	NO	NO	NO	NO	NO	YES
Key Freshness	NO	NO	NO	NO	NO	NO	YES
MITM Attack	NO	YES	YES	NO	YES	YES	YES
Replay Attack	YES	YES	YES	YES	YES	YES	YES
Eavesdropping Attack	NO	NO	NO	NO	NO	YES	YES
Stolen Personal Device	NO	NO	NO	NO	NO	NO	YES
Healthcare Phase	NO	NO	NO	NO	NO	NO	YES

B. Communication cost

During the login and authentication step, the cost of sent messages is calculated. we assumed the identity size is 32 bit, the hash value's size is 160 bits[41], the cipher text value size is 128 bit, we also compare our proposed scheme with other related works based on Table 5 below.

TABLE 5
COMMUNICATION COST

Authors	No of bits	No of messages
Zhang et al.[29]	1568	4
Kaul et al. [30]	768	4
Yassin et al. [39]	576	2
Taher et al.[40]	1660	3
Our Scheme	608	3

VII. CONCLUSIONS

EHRs enable authorized health stakeholders to share structured medical data in order to improve healthcare delivery quality. In these systems, privacy and security are critical, since if sensitive information is leaked, the patient might face serious consequences. Concerns about security and privacy are seen as important barriers in the healthcare system. Remote user authentication is a crucial step in authenticating a person's identity. There have been a number of techniques for remote user authentication, each with their own set of advantages and disadvantages. We propose a secure user authentication scheme for patients in the healthcare system that uses Scyther, a formal security tool, to confirm the security of the proposed work. In the healthcare setting, our work enables the creation of an environment capable of setting, registering, storing, finding, analyzing, authenticating, and validating electronic healthcare information in order to secure patient information. The suggested work uses a lightweight crypto-hash function for the creation of One-Time Passwords (OTP) and symmetric key encryption (Counter mode (CTR mode) to provide good security while maintaining adequate performance. As a result, the primary goal of this study is to develop a reliable

authentication method based on cryptosystem tools in order to address the problems identified in the previous study. Because security threats and demands differ, our solutions focus on patient authentication and privacy. Keep this information private and away from unwanted access. The proposed system will be able to fight off attacks like Man-in-the-Middle, Insider, Replay, and more. It's safe to use features like mutual authentication, anomalies, key management, and other things that are safe. Achieve a balance between speed and security. In the future, we will focus on administrators using two-factor authentication. First-factor biometrics (password and user name) and second-factor biometrics (fingerprint) improve the security of electronic health records patient HER, which is more secure than traditional authentication factors.

CONFLICT OF INTEREST

The authors have no conflict of relevant interest to this article.

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